

# On the Incompatibility of Efficiency and Strategyproofness in Randomized Social Choice

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## Randomized Social Choice

A **social decision scheme (SDS)** maps a preference profile to a **lottery** (probability distribution) over alternatives. Two desirable properties of SDSs are

- **Efficiency:** no agent can be made better off without making another agent worse off, and
- **Strategyproofness:** no agent can benefit from misrepresenting his preferences.

## Related Work

**Gibbard (1973) & Satterthwaite (1975):** Every single-valued, Pareto-optimal, and strategyproof social choice function is a dictatorship.

**Gibbard & Sonnenschein (1977):** Every Pareto optimal and strongly *SD*-strategyproof SDS is a probability mixture of dictatorships.

**Bogomolnaia & Moulin (2001):** There is no anonymous, neutral, *SD*-efficient, and strongly *SD*-strategyproof SDS.

## Preliminaries

Each **agent**  $i$  has a complete and transitive **preference relation**  $R_i$  over a set of **alternatives**  $A$ . A **social decision scheme (SDS)** maps a preference profile  $(R_1, \dots, R_n)$  to a **lottery**  $p$  over  $A$ . An SDS is

- **anonymous** if it is symmetric with respect to agents.
- **neutral** if it is symmetric with respect to alternatives.
- **pairwise** if the outcome depends on the pairwise majority comparisons only.

## Lottery Extensions

Many reasonable lottery extensions can be defined.

**Bilinear dominance (BD):**

$p R_i^{BD} q$  if for all  $x, y$ , where  $x P_i y$ ,  $p(x)q(y) \geq p(y)q(x)$ .

**Stochastic dominance (SD):**

$p R_i^{SD} q$  if for all  $x$ ,  $\sum_{y R_i x} p(y) \geq \sum_{y R_i x} q(y)$ .

**Pairwise comparison (PC):**

$p R_i^{PC} q$  if for all  $x, y$ ,  $\sum_{x P_i y} p(x)q(y) \geq \sum_{x P_i y} q(x)p(y)$ .

**Downward/Upward lexicographic (DL/UL):**

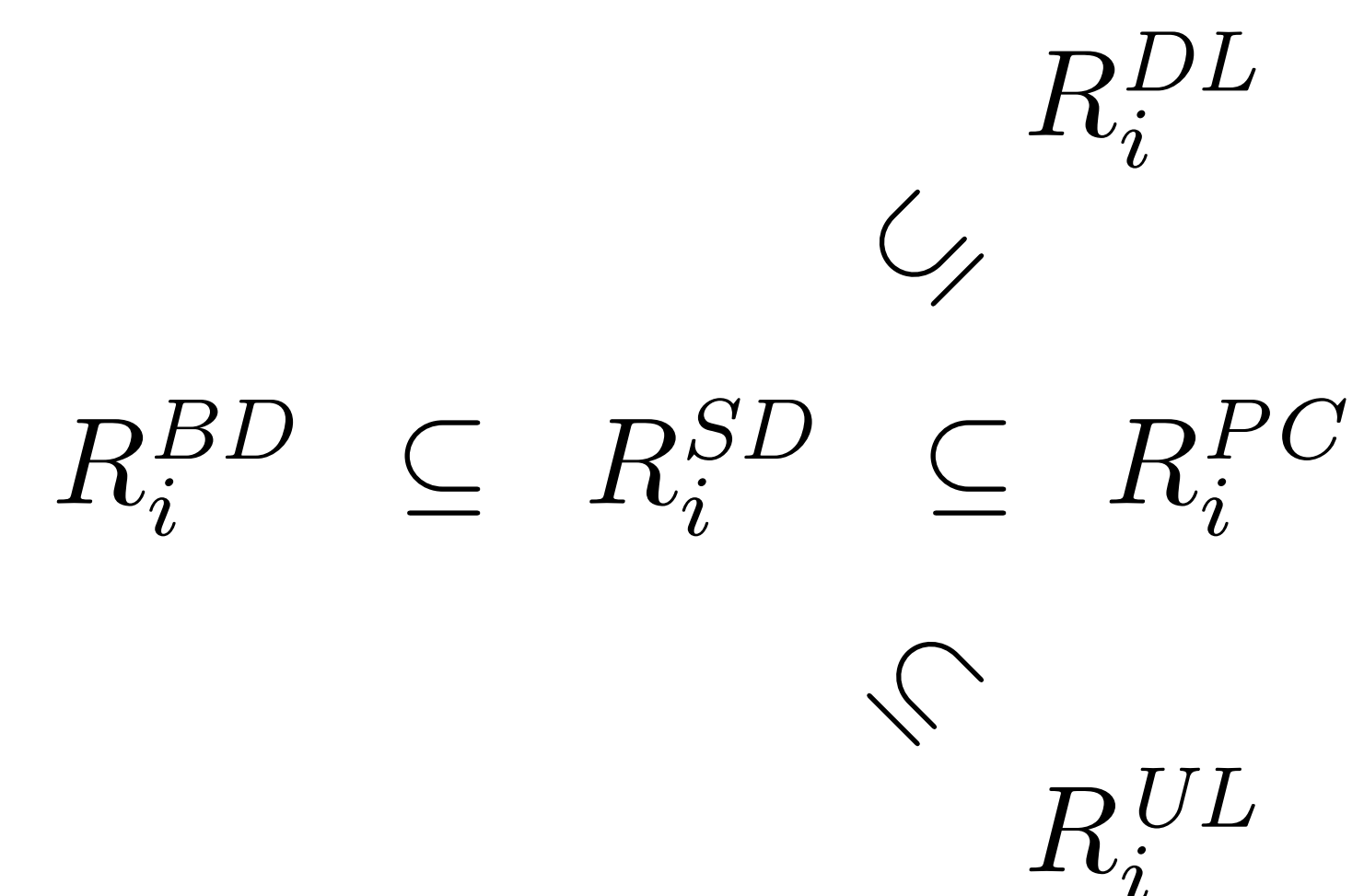
$p R_i^{DL} q$  if there is  $x$  such that  $p(x) > q(x)$  and for all  $y P_i x$ ,  $p(y) = q(y)$ .

$p R_i^{UL} q$  if there is  $x$  such that  $p(x) < q(x)$  and for all  $x P_i y$ ,  $p(y) = q(y)$ .

Example:  $a P_i b P_i c$ ;

$p = \frac{2}{3}a + \frac{1}{3}c$ ;  $q = b$ ; and  $r = \frac{2}{3}b + \frac{1}{3}c$ .

Then  $p P_i^{PC} q$ ;  $p P_i^{DL} q$ ;  $p P_i^{SD} r$ ;  $q P_i^{BD} r$ ; and  $q P_i^{UL} p$ .



## Efficiency and Strategyproofness

Let  $E \in \{BD, SD, PC, DL, UL\}$  be a lottery extension.

- In a preference profile  $R$ , a lottery  $p$  is  **$E$ -efficient** if there is no  $q$  such that  $q R_i^E p$  for all  $i$  and  $q P_i^E p$  for some  $i$ . An SDS  $f$  is  $E$ -efficient if  $f(R)$  is  $E$ -efficient for every preference profile  $R$ .
- An SDS  $f$  is **Pareto-optimal** if  $f(R)$  does only support Pareto-optimal alternatives for every  $R$ .
- An SDS  $f$  is  **$E$ -strategyproof** if there are no  $R, R', i$  such that  $R_j = R'_j$  for all  $j \neq i$  and  $f(R') P_i^E f(R)$ .
- An SDS  $f$  is **strongly  $E$ -strategyproof** if  $f(R) R_i^E f(R')$ , whenever  $R_j = R'_j$  for all  $j \neq i$ .

## Impossibility Theorems

- There is no anonymous, neutral, *PC*-efficient, and *PC*-strategyproof SDS.
- There is no anonymous, *UL*-efficient, and *UL*-strategyproof SDS.\*
- There is no pairwise, Pareto-optimal, and *BD*-strategyproof SDS.\*\*
- There is no anonymous, neutral, Pareto-optimal, and *BD*-group-strategyproof SDS.\*\*\*

**Conjecture:** There is no anonymous, *SD*-efficient, and *SD*-strategyproof SDS.

Anonymity is required for all our results as **serial dictatorship** satisfies all considered notions of efficiency and strategyproofness.

\*Interestingly, random dictatorship satisfies anonymity, neutrality, *DL*-efficiency, and *DL*-strategyproofness.

\*\*This strengthens a result by Aziz, Brandt, and Brill (2013).

\*\*\*This strengthens a result by Bogomolnaia, Moulin, and Stong (2005).